

Persistent Homology and Nested Dissection

Don Sheehy
University of Connecticut

joint work with Michael Kerber and Primoz Skraba

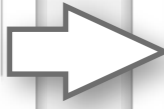
A Topological Data Analysis Pipeline

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(Lipschitz)

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(to approx. the function)

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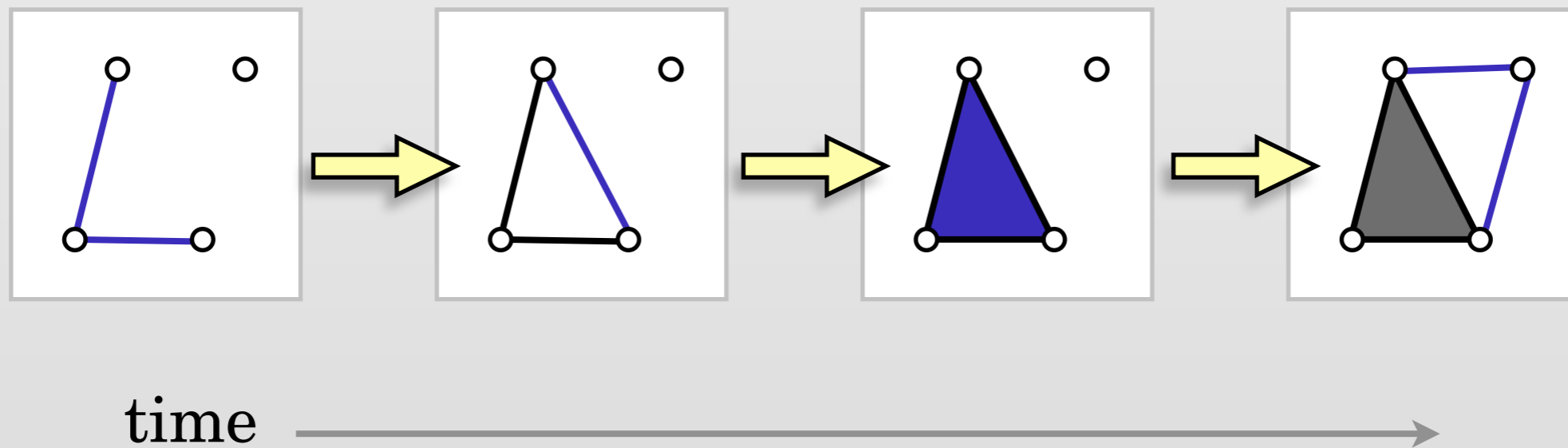
How to compute it?

Q: Can we use the geometry to speed up the persistence computation?

We want to build a **filtered simplicial complex**.

Associate a **birth time** with each simplex in complex K .

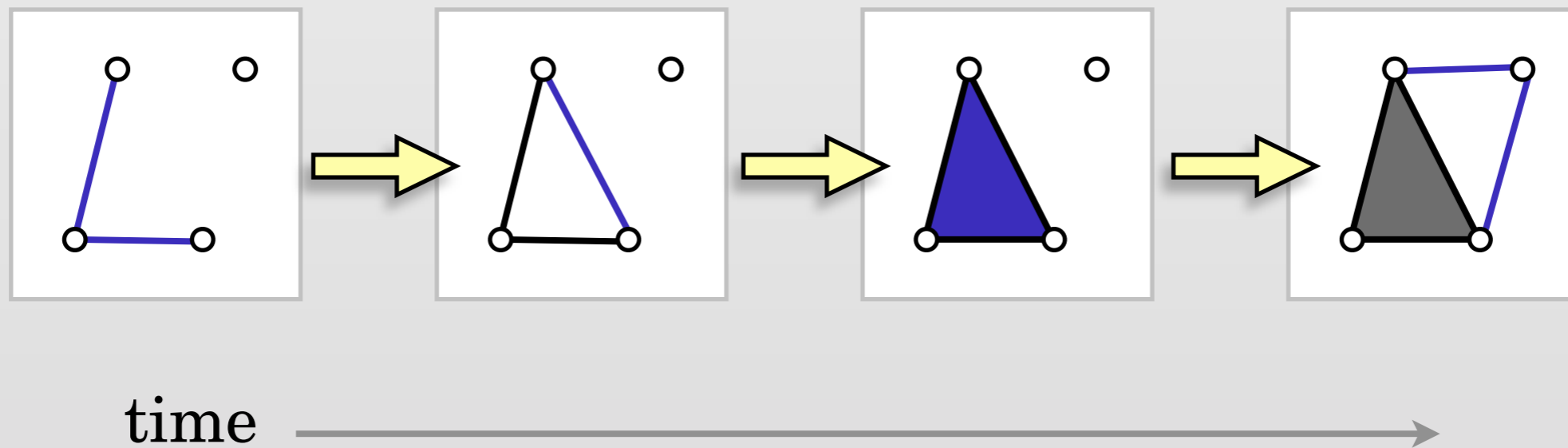
At time α , we have a complex K_α consisting of all simplices born at or before time α .



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This can be represented by its boundary matrix D , with a fixed row order.

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It's just Gaussian elimination!

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Clear and Compress [BKR13]

Nested Dissection

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A method for solving symmetric positive definite linear systems.

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Inverting A can be done in $O(n^{\beta\omega})$ time.
Also works for computing ranks of singular,
nonsymmetric matrices over finite fields.

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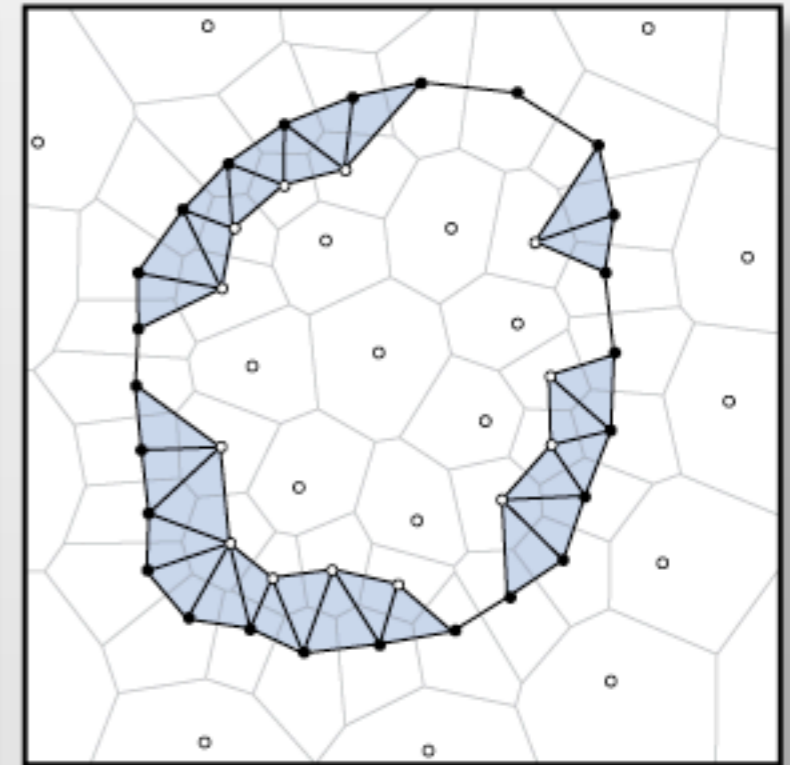
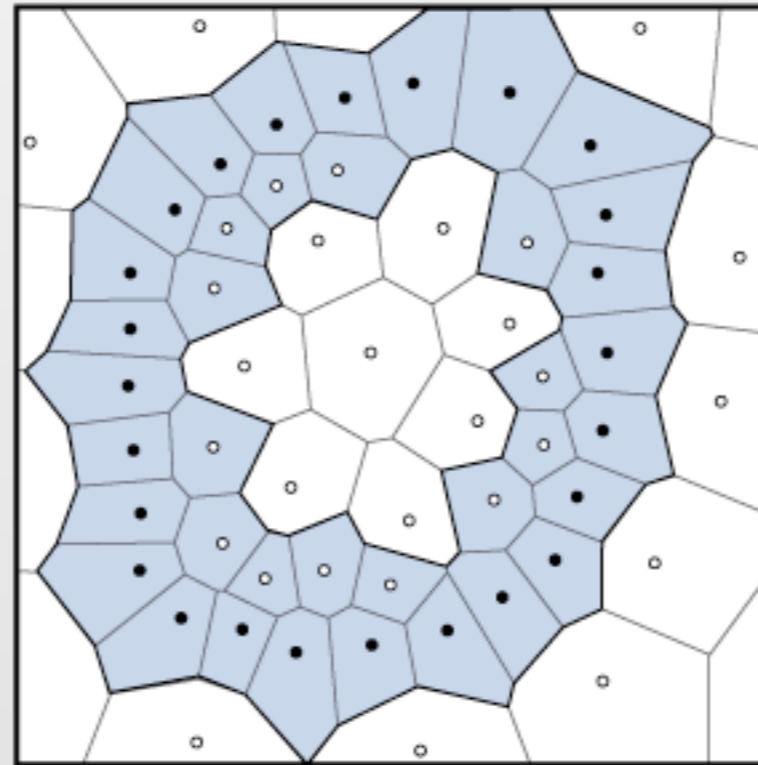
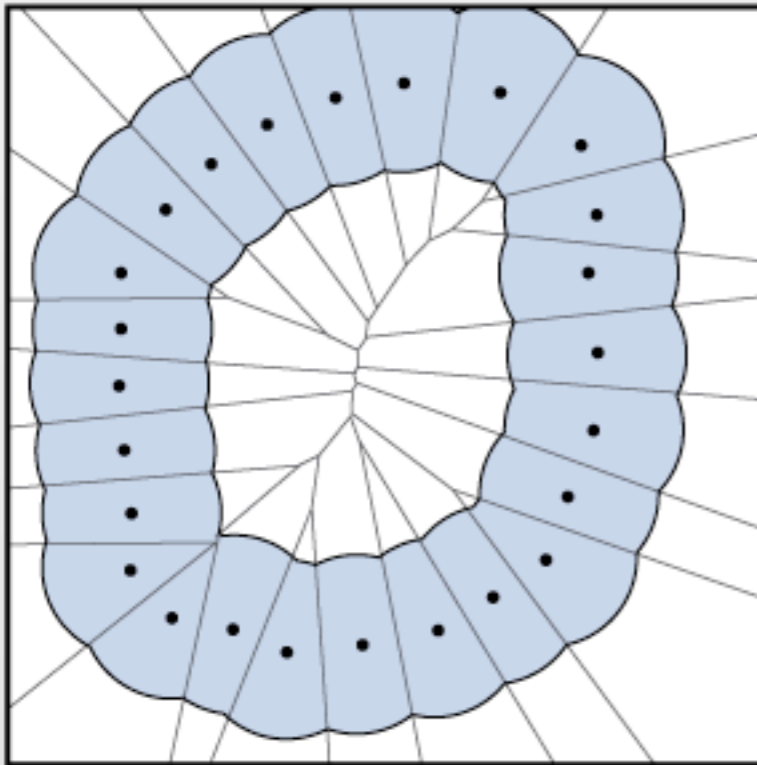
The other trick:

The geometric separators of Miller-Teng-Thurston-Vavasis can also give separators for minors of the boundary matrix of a quality mesh.

Mesh Filtrations

Geometric
Approximation

Topologically
Equivalent

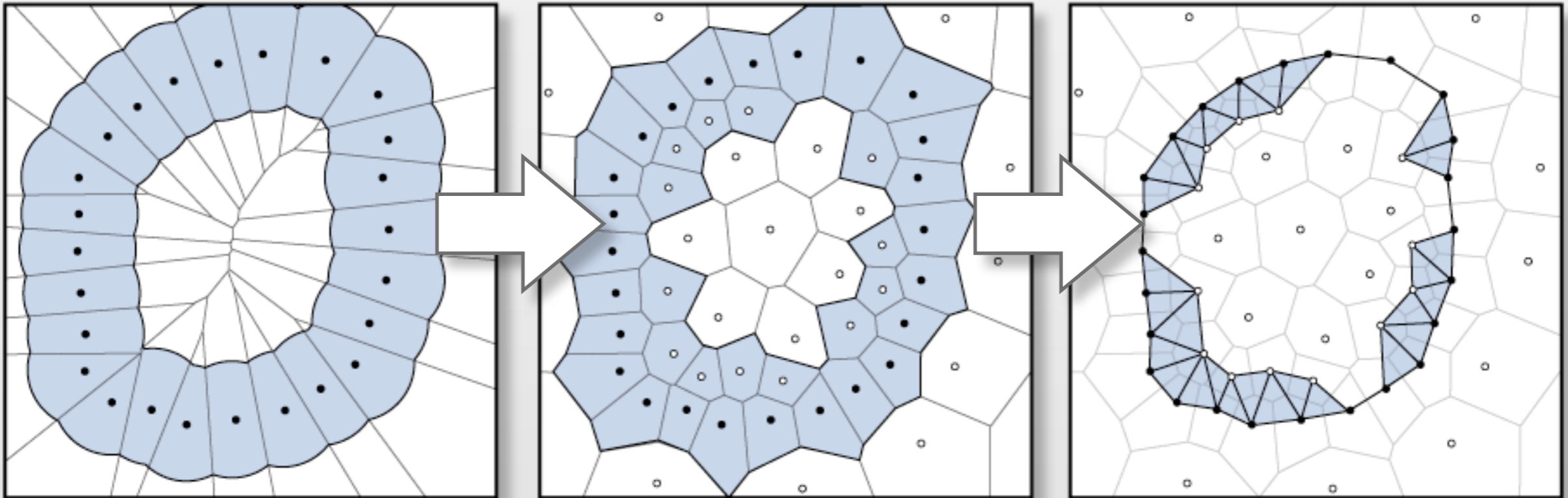


1. Compute the function on the vertices.
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3. Filter the Delaunay triangulation appropriately.

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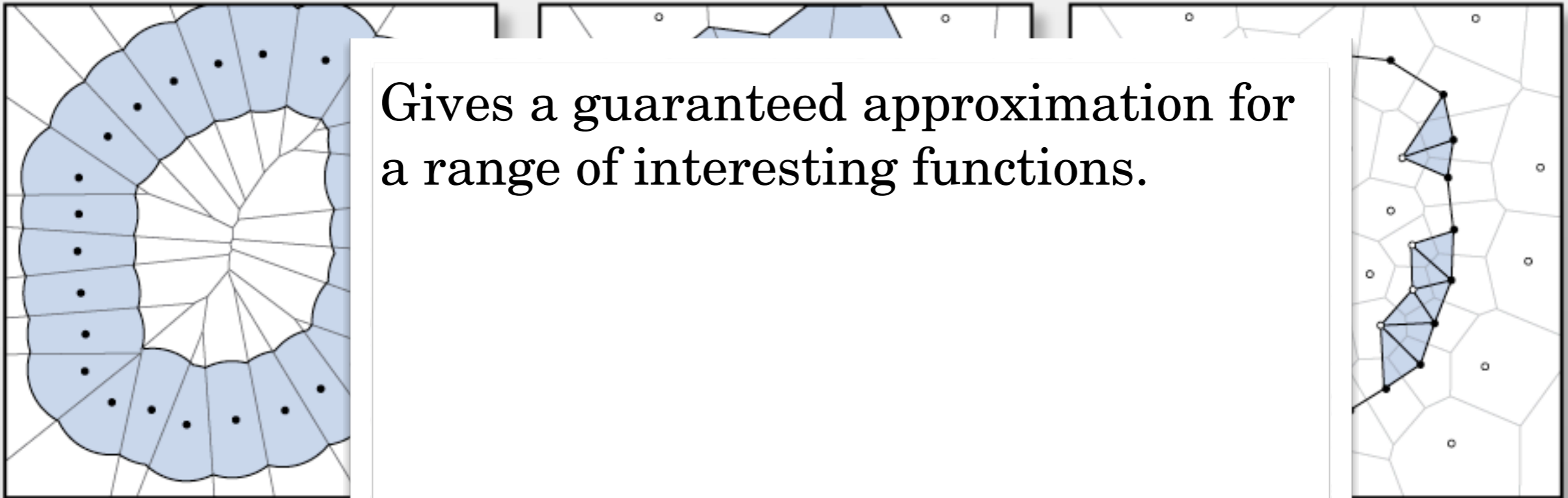


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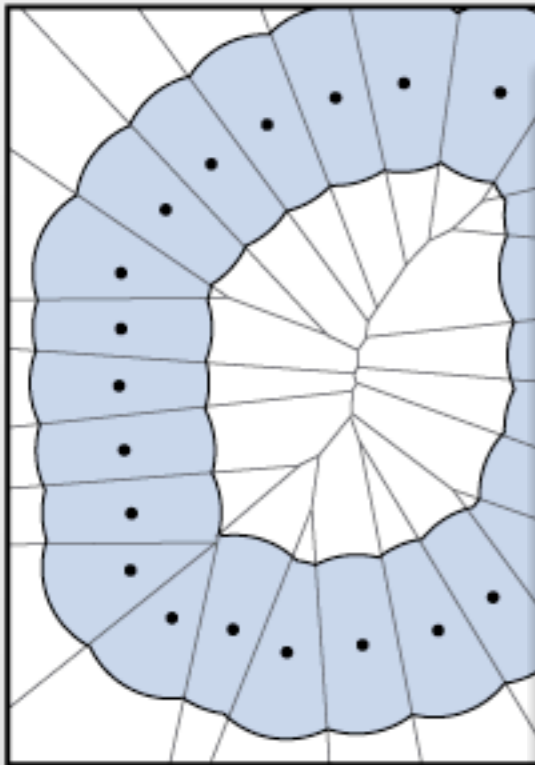
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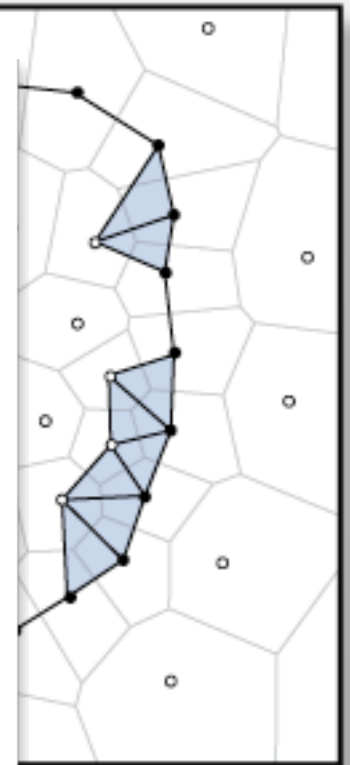
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Bonus: The theory of **geometric separators** was invented for graphs coming from meshes like this.

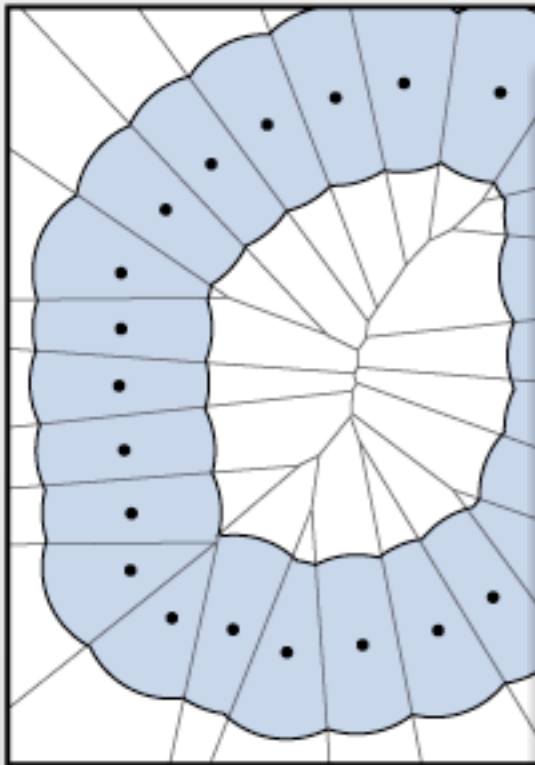


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We extend the geometric separators to complex separators.

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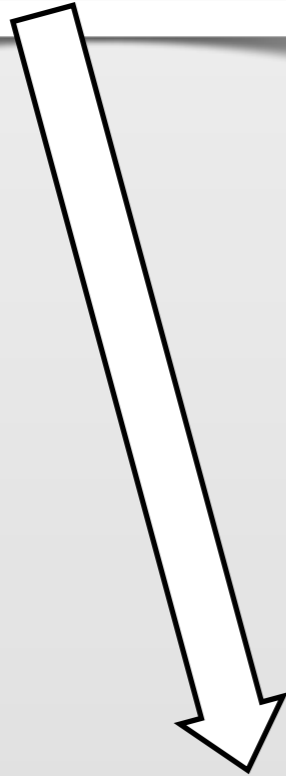
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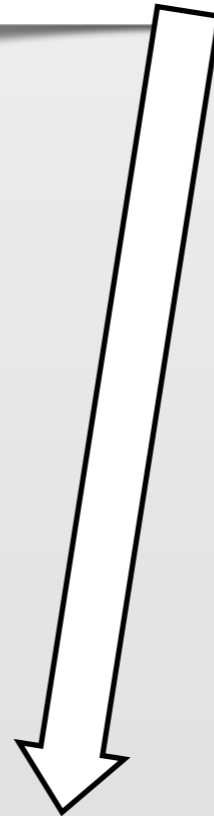
Output-Sensitive
Persistence Algorithm

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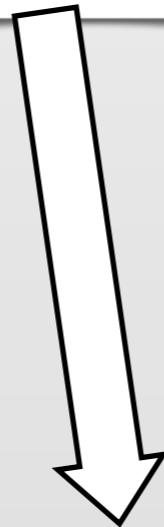
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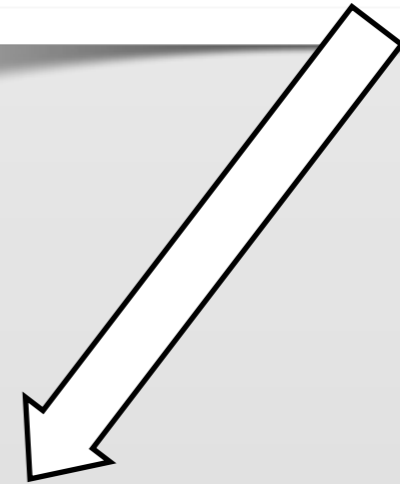
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Theorem.

Let F be a filtration on the Delaunay triangulation of a set of τ -well-spaced points. For a constant $\Gamma \geq 0$, the subset D_Γ of the persistence diagram of F consisting of those pairs with persistence at least Γ can be computed in $O(|D_{(1-\delta)\Gamma}|n^{\omega(1-\frac{1}{d})})$ time, where $D_{(1-\delta)\Gamma}$ is the set of pairs in the persistence diagram with persistence at least $(1-\delta)\Gamma$ for any constant $\delta > 0$.

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Thanks.

